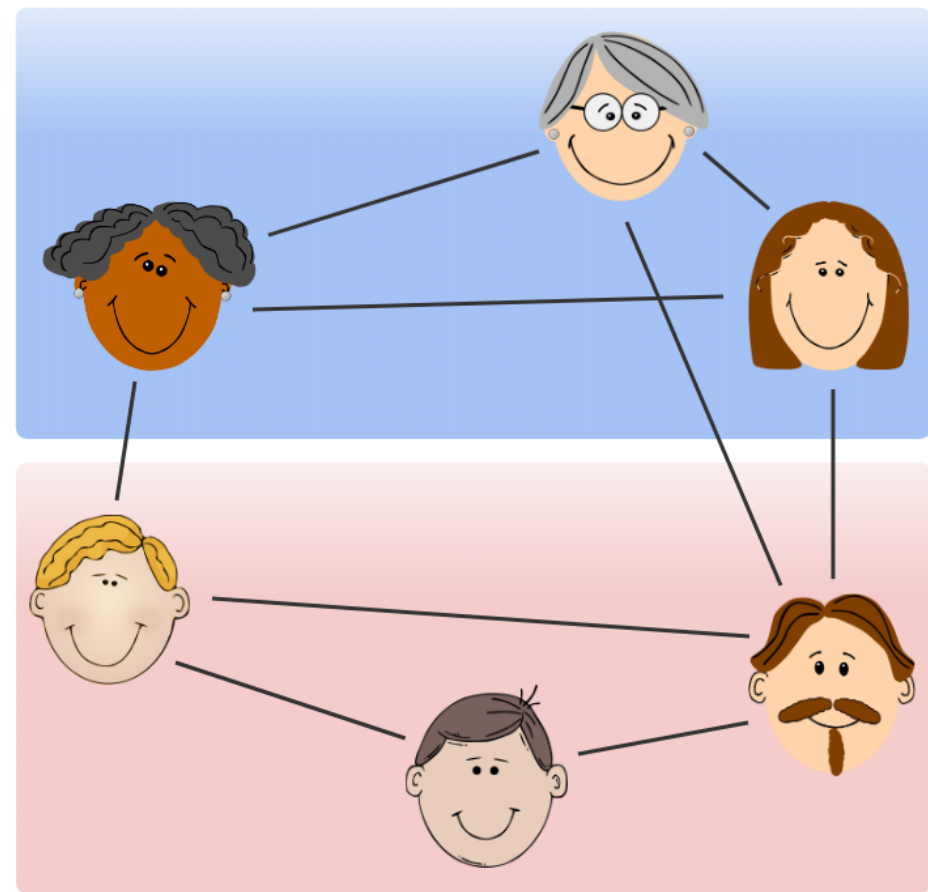


Spectral clustering (SC)

SC is the method of choice for clustering the nodes of a graph.

Example: want to find clusters in a friendship network — many connections within clusters, few connections between clusters



Friendship network and a clustering into two clusters obtained from running SC. According to the fairness notion of [Chierichetti et al. \(2017\)](#) this clustering is highly unfair with respect to the two demographic groups of males and females (see below).

There are several versions of SC ([von Luxburg, 2007](#)); here we focus on *unnormalized SC*, but our results hold similarly for *normalized SC* too.

Formally:

$G = (V, E)$... undirected graph on $V = [n]$; $W \in \mathbb{R}^{n \times n}$... (weighted) adjacency matrix with $W_{ij} > 0$ iff there is an edge between i and j ; k ... number of clusters

SC aims to partition V into k clusters with minimum value of the *RatioCut objective function* defined as follows: for a clustering $V = C_1 \dot{\cup} \dots \dot{\cup} C_k$ it is

$$\text{RatioCut}(C_1, \dots, C_k) = \sum_{l=1}^k \frac{\text{Cut}(C_l, V \setminus C_l)}{|C_l|} \quad \text{with} \quad \text{Cut}(C_l, V \setminus C_l) = \sum_{\substack{i \in C_l, \\ j \in V \setminus C_l}} W_{ij}.$$

Let D be the degree matrix (with $D_{ii} = \sum_{j \in [n]} W_{ij}$) and $L = D - W$ be the *graph Laplacian matrix*. Encoding a clustering $V = C_1 \dot{\cup} \dots \dot{\cup} C_k$ by $H \in \mathbb{R}^{n \times k}$ with

$$H_{il} = \begin{cases} 1/\sqrt{|C_l|}, & i \in C_l, \\ 0, & i \notin C_l, \end{cases} \quad (1)$$

we have $\text{RatioCut}(C_1, \dots, C_k) = \text{Tr}(H^T L H)$. SC relaxes the exact problem

$$\min_{H \in \mathbb{R}^{n \times k}} \text{Tr}(H^T L H) \quad \text{subject to } H \text{ is of form (1)} \quad (2)$$

and solves

$$\min_{H \in \mathbb{R}^{n \times k}} \text{Tr}(H^T L H) \quad \text{subject to } H^T H = I_k \quad (3)$$

instead. A solution to (3) is given by a matrix H that contains some orthonormal eigenvectors corresponding to the k smallest eigenvalues of L as columns. Then, a clustering of V is inferred from H by applying k -means clustering to the rows of H .

Spectral clustering with fairness constraints

Let V contain h demographic groups V_s such that $V = \dot{\cup}_{s \in [h]} V_s$. [Chierichetti et al. \(2017\)](#) proposed a *notion of fairness* for clustering: **in every cluster, each group V_s should be represented with (approximately) the same fraction as in the whole data set V .**

We can incorporate this goal into (2) via a linear constraint on H :

The face images were found on <https://openclipart.org> and are in the public domain.

Algorithm 1 Unnormalized SC with fairness constraints

Input: adj. matrix $W \in \mathbb{R}^{n \times n}$; $k \in \mathbb{N}$; group-membership vectors $f^{(s)} \in \{0, 1\}^n$

Output: a clustering of $[n]$ into k clusters

- compute the Laplacian matrix $L = D - W$
- let F be a matrix with columns $f^{(s)} - \frac{|V_s|}{n} \cdot \mathbf{1}_n$, $s \in [h-1]$
- let the columns of Z form an orthonormal basis of the nullspace of F^T
- compute the k smallest (respecting multiplicities) eigenvalues of $Z^T L Z$ and the corresponding orthonormal eigenvectors (written as columns of Y)
- apply k -means clustering to the rows of $H = ZY$

Lemma 1 For $s \in [h]$, let $f^{(s)} \in \{0, 1\}^n$ be the group-membership vector of V_s , that is $f_i^{(s)} = 1$ if $i \in V_s$ and $f_i^{(s)} = 0$ otherwise. Let $V = C_1 \dot{\cup} \dots \dot{\cup} C_k$ be a clustering that is encoded as in (1). We have, for every $l \in [k]$,

$$\forall s \in [h-1] : \sum_{i=1}^n \left(f_i^{(s)} - \frac{|V_s|}{n} \right) H_{il} = 0 \quad \Leftrightarrow \quad \forall s \in [h] : \frac{|V_s \cap C_l|}{|C_l|} = \frac{|V_s|}{n}. \quad (6)$$

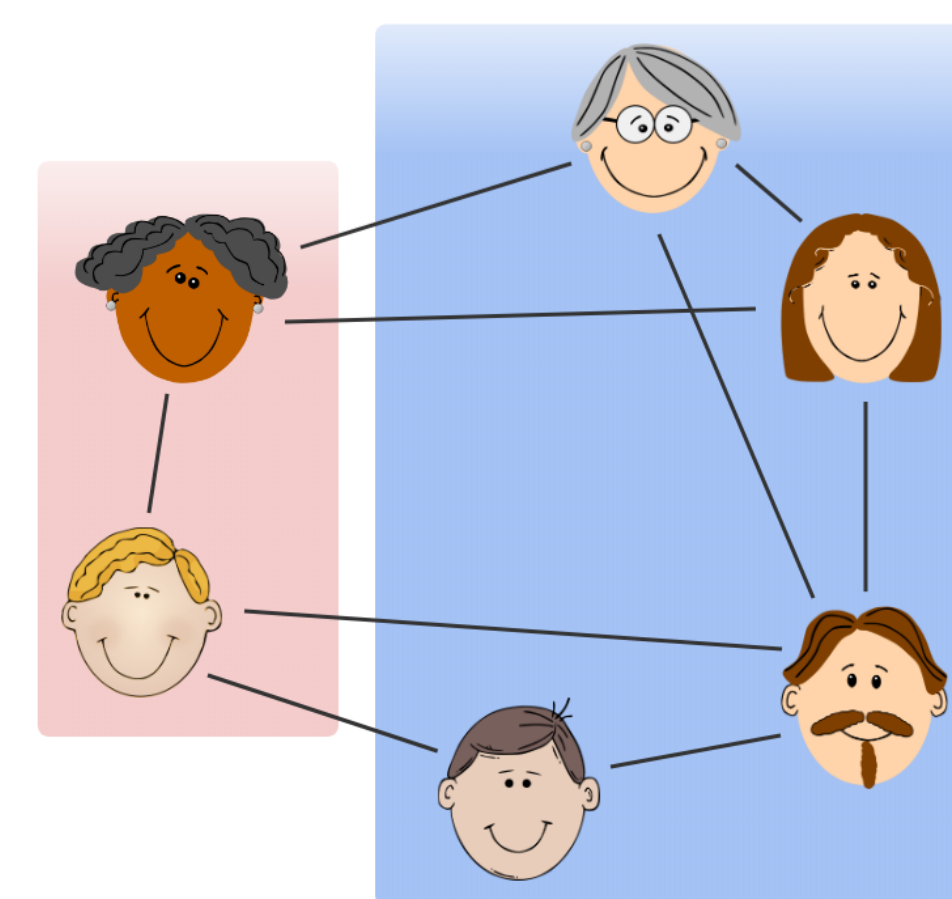
Let $F \in \mathbb{R}^{n \times (h-1)}$ comprise the vectors $f^{(s)} - (|V_s|/n) \cdot \mathbf{1}_n$, $s \in [h-1]$, as columns. According to Lemma 1, instead of (3) we should solve

$$\min_{H \in \mathbb{R}^{n \times k}} \text{Tr}(H^T L H) \quad \text{subject to } H^T H = I_k \text{ and } F^T H = 0_{(h-1) \times k}. \quad (4)$$

Our approach is *analogous to existing versions of constrained SC* that try to incorporate must-link constraints into the SC framework (e.g., [Yu and Shi, 2004](#)).

How to solve (4)? Let $Z \in \mathbb{R}^{n \times (n-h+1)}$ be a matrix whose columns form an orthonormal basis of the nullspace of F^T . We can substitute $H = ZY$ for $Y \in \mathbb{R}^{(n-h+1) \times k}$, and then we end up with a problem of the form (3).

Our strategy is summarized in Algorithm 1. Note that, in general, there is no guarantee on how fair the output of Algorithm 1 is. However, often it is quite fair:



A clustering into two clusters obtained from running Algorithm 1. According to the fairness notion of [Chierichetti et al. \(2017\)](#) this clustering is perfectly fair.

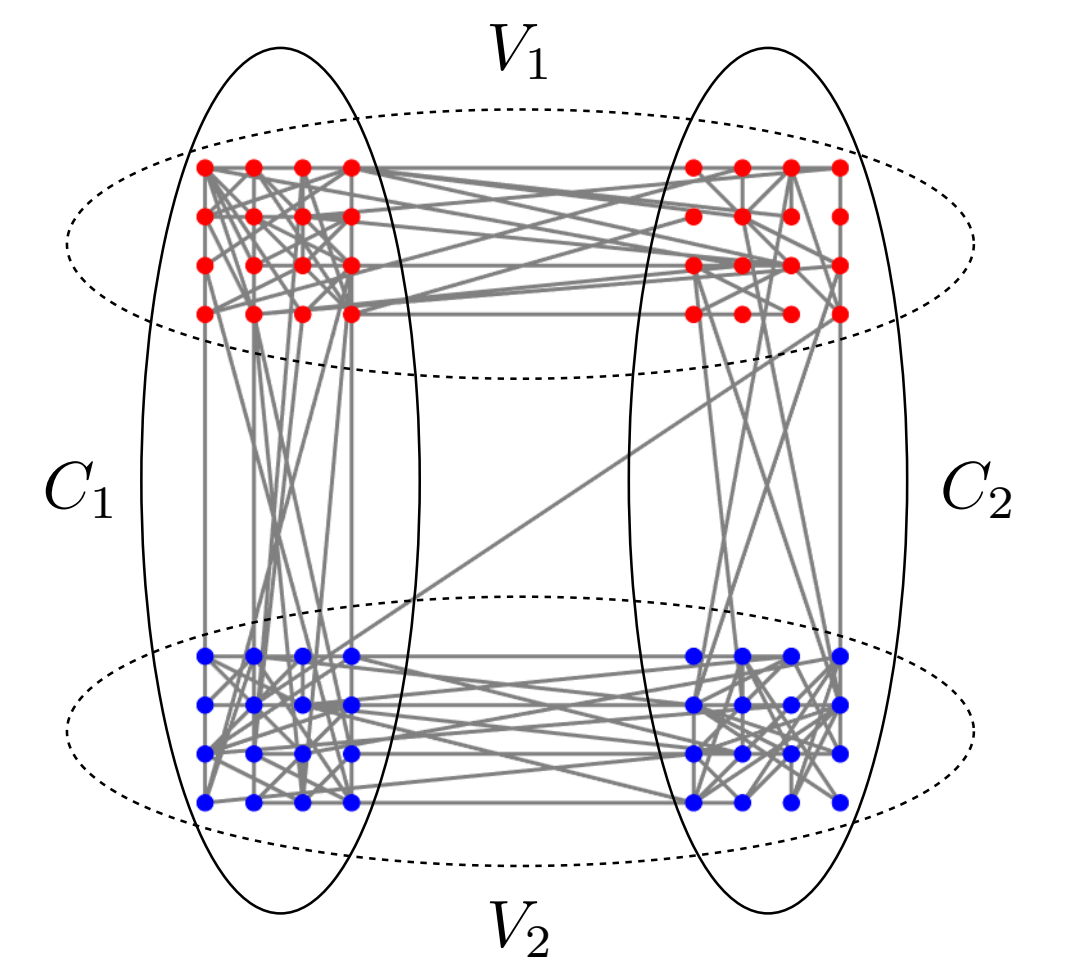
Analysis on variant of the stochastic block model (SBM)

We analyze our fair version of SC on a natural variant of the famous SBM ([Holland et al., 1983](#)): let $V = [n]$ comprise h groups $V = \dot{\cup}_{s \in [h]} V_s$ and be partitioned into k ground-truth clusters $V = C_1 \dot{\cup} \dots \dot{\cup} C_k$ such that $|V_s \cap C_l|/|C_l| = \eta_s$, $s \in [h]$, $l \in [k]$, for some $\eta_s \in (0, 1)$ (\Rightarrow the ground-truth clustering is perfectly fair). We define a random graph on V by connecting vertices i and j with probability $\Pr(i, j)$, where

$$\Pr(i, j) = \begin{cases} a, & i \text{ and } j \text{ in same cluster and in same group,} \\ b, & i \text{ and } j \text{ not in same cluster, but in same group,} \\ c, & i \text{ and } j \text{ in same cluster, but not in same group,} \\ d, & i \text{ and } j \text{ not in same cluster and not in same group} \end{cases} \quad (5)$$

for some $a > b > c > d$. An example of such a graph can be seen in the next figure.

Example of a graph generated from our variant of the SBM. There are two meaningful clusterings into two clusters: $V = C_1 \dot{\cup} C_2$ and $V = V_1 \dot{\cup} V_2$. However, only the first one is fair.



For such a graph, standard SC is likely to return the unfair clustering $V = V_1 \dot{\cup} V_2$ while Algorithm 1 returns the fair clustering $V = C_1 \dot{\cup} C_2$ with high probability:

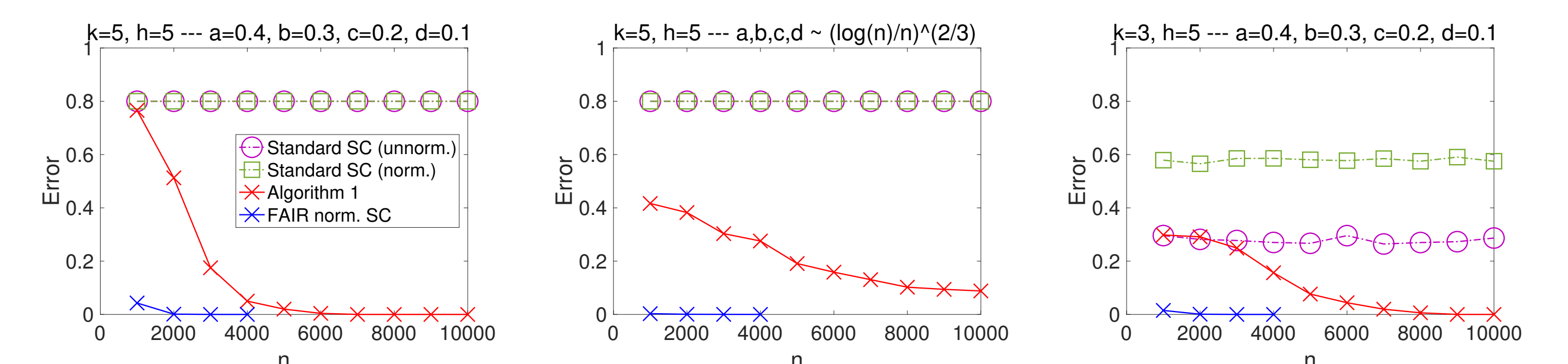
Theorem 1 Let $V = [n]$ comprise h groups $V = \dot{\cup}_{s \in [h]} V_s$ and be partitioned into k ground-truth clusters $V = C_1 \dot{\cup} \dots \dot{\cup} C_k$ such that

$$|V_s| = \frac{n}{h}, \quad |C_l| = \frac{n}{k}, \quad \frac{|V_s \cap C_l|}{|C_l|} = \frac{1}{h}, \quad s \in [h], l \in [k]. \quad (6)$$

Let G be a random graph constructed according to our variant of the SBM (5) with $a > b > c > d$ and $a \geq C \ln n/n$ for some $C > 0$. Assume that we run Algorithm 1 on G , where we apply a $(1+M)$ -approximation algorithm to the k -means problem encountered in the last step, for some $M > 0$. Then there exist $\tilde{C}, \tilde{C} > 0$ such that the following is true: if $a(c-d)^{-2} k^3 n^{-1} \ln n < \tilde{C}(1+M)^{-1}$, then with probability at least $1 - n^{-1}$, the clustering returned by Algorithm 1 “misclassifies” at most $\tilde{C} a(c-d)^{-2} k^2 (1+M) \ln n$ many vertices.

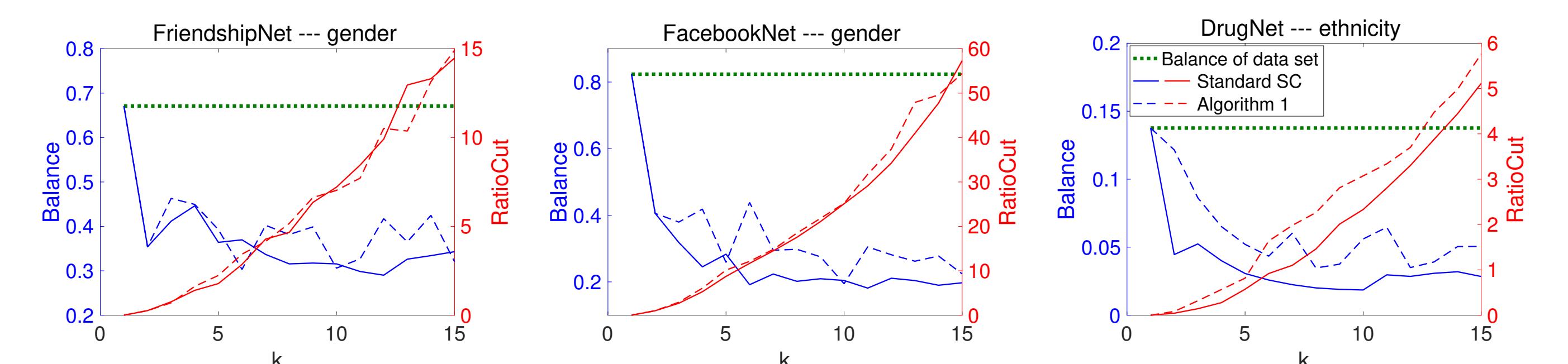
Experiments

SBM



Error (fraction of “misclassified” vertices) as a function of n . In the third plot, C_1, C_2 and C_3 have different sizes and Assumption (6) is not satisfied.

Real networks



Average balance of clusters and RatioCut value as a function of k .

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